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Christian Servin

*El Paso Community College*, cservin@gmail.com

Gerardo Muela

*University of Texas at El Paso*, gdmuela@miners.utep.edu

Vladik Kreinovich

*University of Texas at El Paso*, vladik@utep.edu

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Article

# Fuzzy Analogues of Sets and Functions Can Be Uniquely Determined from the Corresponding Ordered Category: A Theorem

Christian Servin <sup>1,†</sup>, Gerardo D. Muela <sup>2,†</sup> and Vladik Kreinovich <sup>2,†</sup>  0000-0002-1244-1650\*

<sup>1</sup> El Paso Community College; cservin@gmail.com

<sup>2</sup> University of Texas at El Paso; gdmuela@miners.utep.edu, vladik@utep.edu

\* Correspondence: vladik@utep.edu; Tel.: +1-915-747-6951

† All the authors contributed equally to this work.

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**Abstract:** In modern mathematics, many concepts and ideas are described in terms of category theory. From this viewpoint, it is desirable to analyze what can be determined if, instead of the basic category of sets, we consider a similar category of fuzzy sets. In this paper, we describe a natural fuzzy analog of the category of sets and functions, and we show that, in this category, fuzzy relations (a natural fuzzy analogue of functions) can be determined in category terms – of course, modulo 1-1 mapping of the corresponding universe of discourse and 1-1 re-scaling of fuzzy degrees.

**Keywords:** fuzzy set; ordered category; category of fuzzy sets

## 1. Introduction

**What are categories: a brief reminder.** While set theory remains the foundations for mathematics, in modern mathematics, many concepts and ideas are described in terms of category theory.

A *category* is a tuple  $(\text{Ob}, \text{Mor}, :, \text{id}, \circ)$ , where:

- $\text{Ob}$  is the set whose elements are called *objects*,
- $\text{Mor}$  is a set whose elements are called *morphisms*,
- $:$   $\text{Mor} \rightarrow \text{Ob} \times \text{Ob}$  is a mapping that assigns, to each morphism  $f \in \text{Mor}$  a pair of objects  $(a, b) \in \text{Ob} \times \text{Ob}$ ; this is denoted by  $f : a \rightarrow b$ ; the object  $a$  is called  $f$ 's *domain*, and  $b$  is called  $f$ 's *range*;
- $\text{id}$  is a mapping that assigns, to each object  $a \in \text{Ob}$ , a morphism  $\text{id}_a : a \rightarrow a$ ; and
- $\circ$  is a mapping that assigns, to each pair of morphisms  $f : a \rightarrow b$  and  $g : b \rightarrow c$  for which the range of  $f$  is equal to the domain of  $g$ , a new morphism  $g \circ f : a \rightarrow c$  so that for every  $f : a \rightarrow b$ , we have  $\text{id}_b \circ f = f \circ \text{id}_a = f$ .

For example:

- We can have a category *Set* in which objects are sets and morphisms are functions.
- We can have a category *Top* in which objects are topological spaces and morphisms are continuous mappings.
- We can have a category *Lin*, in which objects are linear spaces, and morphisms are linear mappings, etc.

Many mathematical concepts can be reformulated in terms of an appropriate category.

28 **What happens in the fuzzy case?** If we allow fuzzy sets (see, e.g., [2–6]), what is a natural  
 29 analog of the category Set? In the category Set, morphisms from  $a$  to  $b$  are functions. In the crisp case,  
 30 for each function  $f : a \rightarrow b$  and for each element  $x \in a$ , we have a unique value of  $y = f(x) \in b$ .

31 Fuzzy means that for each  $x \in a$ , instead of a single value  $y = f(x) \in b$ , we may have different  
 32 possible values  $y \in b$ , with different degrees of confidence. In general, we can have all possible values  
 33  $y \in b$ . For each  $x \in a$  and for each  $y \in b$ , we have a degree  $R_f(x, y)$  to which  $y$  is a possible value of  
 34  $f(x)$ . Thus, a natural fuzzy analog of a function is a fuzzy relation.

Composition  $g \circ f$  of fuzzy relations  $f : a \rightarrow b$  and  $g : b \rightarrow c$  can be defined in the usual way.  
 Namely, we want to know, for each pair of elements  $x \in a$  and  $c \in c$ , to what extent there exists a  $y \in b$   
 for which  $f$  brings us from  $a$  to  $b$  and  $g$  brings us from  $y$  to  $c$ . If we interpret “and” as min and there  
 exists (an infinite “or”) as max, then the above description translates into the following formula:

$$R_{g \circ f}(x, z) = \max_y \min(R_f(x, y), R_g(y, z)). \quad (1)$$

35 Since we have fuzzy relations, there is no need to explicitly describe the domain of each morphism:  
 36 if for some  $x \notin a$ , the value  $f(x)$  is not defined, this simply means that for this  $x$ , we have  $R_f(x, y) = 0$   
 37 for all  $y \in b$ . Similarly, there is no need to describe the range,

38 Thus, without losing generality, we can assume that the relation  $R_f(x, y)$  is defined for all  $x \in U$   
 39 and  $y \in U$ .

40 Thus, without losing generality, we can assume that we have only one object – the universal set  $U$ .  
 41 Morphisms are then fuzzy relations, with the usual composition relation (1).

42 **Need for an ordered category.** In the crisp case, every property is either true or false.

43 As we gain more information, we may get more confident in our knowledge. For example, we  
 44 may start with the situation in which, for a given  $x$ , several different values  $f(x)$  are possible, but after  
 45 acquiring new information, we are becoming more and confident that there is only one possible value  
 46  $y_0$  of  $f(x)$ . This means that for the remaining value  $y_0$ , the degree of possibility  $R_f(x, y_0)$  remains the  
 47 same, but for all  $y \neq y_0$ , the corresponding degree  $R_f(x, y)$  decreases. To capture this phenomenon, it  
 48 is reasonable to supplement the category structure with the corresponding component-wise ordering  
 49 between fuzzy relations (morphisms):  $f \leq f'$  if and only if  $R_f(x, y) \leq R_{f'}(x, y)$  for all  $x$  and  $y$ .

50 **Formulation of the problem.** What can be defined based on this category-theory formulation? Can  
 51 we uniquely determine the elements of the Universe of discourse  $U$  and the corresponding relations  
 52 based on the categorical information?

## 53 2. Results

54 **Towards a precise formulation of the problem.** It is easy to see that if we have a 1-1 mapping  
 55  $\pi : U \rightarrow U$  of the Universe of discourse  $U$  onto itself (i.e., a bijection), then the corresponding  
 56 mapping, then the corresponding transformation  $R(x, y) \rightarrow R(\pi(x), \pi(y))$  is an *automorphism* of the  
 57 corresponding category in the sense that it preserves the identity, composition, and order.

58 Similarly, if we have a 1-1 monotonic mapping  $H : [0, 1] \rightarrow [0, 1]$ , then the transformation  
 59  $R(x, y) \rightarrow H(R(x, y))$  is also such an automorphism. Indeed, since we only consider order  
 60 between degrees, monotonic transformation of degrees should not change anything.

61 It turns out that modulo this simple equivalence, we can uniquely determine all the elements  $x \in$   
 62  $U$  and all the relations  $R(x, y)$  from the ordered category, i.e., in precise terms, that every automorphism  
 63 is a composition of the automorphisms of the above two types. The proof of this result will be based  
 64 on an explicit description of elements of  $U$  and relations  $R_f(x, y)$  in category terms.

65 Let us describe the problem in precise terms.

66 **Definition 1.** By an ordered category, we mean a category in which for every two objects  $a$  and  $b$ , there is a  
 67 partial order  $\leq$  on the set  $\text{Mor}(a, b)$  of all morphisms from  $a$  to  $b$ .

68 **Definition 2.** Let  $U$  be a set; we will call it the Universe of discourse. By a  $U$ -fuzzy ordered category, we  
69 mean an ordered category in which:

- 70 • the only object is the set  $U$ ,
- 71 • morphisms are fuzzy relations, i.e., mappings  $R : U \times U \rightarrow [0, 1]$ ,
- 72 • the morphism  $\text{id}$  is defined as the mapping for which  $\text{id}(x, x) = 1$  and  $\text{id}(x, y) = 0$  for  $x \neq y$ ,
- the composition of morphisms is defined by the formula

$$(g \circ f)(x, z) = \max_y \min(f(x, y), g(y, z)),$$

73 and

- 74 • the order between the morphisms is the componentwise order:  $f \leq g$  means that  $f(x, y) \leq g(x, y)$  for all  
75  $x$  and  $y$ .

76 The  $U$ -fuzzy ordered category will be denoted by  $F_U$ .

77 *Comment.* One can easily see that this is indeed a category, i.e., that the composition of morphisms  
78 is associative, and the composition of any morphism  $f$  with the identity morphism  $\text{id}$  is equal to  $f$ :  
79  $f \circ \text{id} = \text{id} \circ f = f$ .

80 **Definition 3.** An automorphism of an ordered category is a pair consisting of bijections  $F : \text{Ob} \rightarrow \text{Ob}$  and  
81  $G : \text{Mor} \rightarrow \text{Mor}$  for which:

- 82 • for all  $f, a$ , and  $b$ , we have  $f : a \rightarrow b$  if and only if  $G(f) : F(a) \rightarrow F(b)$ ;
- 83 • for all  $f$  and  $g$ , we have  $G(f \circ g) = G(f) \circ G(g)$ ,
- 84 • for all  $a$ , we have  $G(\text{id}_a) = \text{id}_{F(a)}$ , and
- 85 • for all  $f$  and  $g$ , we have  $f \leq g$  if and only if  $G(f) \leq G(g)$ .

86 **Proposition.** Let  $\pi : U \rightarrow U$  be a bijection of  $U$ , and let  $H : [0, 1] \rightarrow [0, 1]$  be an increasing bijection of the  
87 interval  $[0, 1]$ . Then, the mapping  $G_{\pi, H}$  that maps each morphism  $f(x, y)$  into a morphism  $(G_{\pi, H}(f))(x, y) =$   
88  $H(\pi(x), \pi(y))$  is an automorphism of the category  $F_U$ .

89 Our main result is that these are the only automorphisms of the category  $F_U$ .

90 **Theorem.** For every set  $U$ , every automorphism of the ordered category  $F_U$  has the form  $G_{\pi, H}$  for  
91 some bijection  $\pi : U \rightarrow U$  and for some monotonic bijection  $H : [0, 1] \rightarrow [0, 1]$ .

92 *Comment.* This may not be very clear from the formulation of the result, but the proof will show that  
93 we can determine elements of the set  $U$  and values of the mappings  $f(x, y)$  in category terms, i.e., we  
94 can indeed define fuzzy relations – a natural fuzzy analogue of functions – in category terms.

### 95 3. Proofs

#### 96 3.1. Proof of the Proposition

97 This proposition is easy to prove: a permutation  $\pi$  does not change anything, and the increasing  
98 bijection does not change the order.

#### 99 3.2. Proof of the Theorem

100  $1^\circ$ . First, we can describe the morphism  $f_0$  for which  $f_0(x, y) = 0$  for all  $x$  and  $y$  in ordered-category  
101 terms, as the only morphism  $f$  for which  $f \leq g$  for all morphisms  $g$ .

102 Indeed, clearly  $f_0 \leq g$  for all  $g$ . Vice versa, if  $f \leq g$  for all  $g$ , then, in particular,  $f \leq f_0$ , i.e.,  
103  $f(x, y) \leq f_0(x, y) = 0$  for all  $x$  and  $y$ , and since  $f(x, y) \in [0, 1]$ , this means that indeed  $f(x, y) = 0$  for  
104 all  $x$  and  $y$ .

105 2°. Let us first characterize all the morphisms  $f \neq f_0$  for which the set  $\{g : g \leq f\}$  is linearly ordered.  
 106 Since an automorphism preserves order, every automorphism maps such morphisms into morphisms  
 107 with the same property.

108 Specifically, we will prove that a morphism has this property if and only if we have  $f(x, y) > 0$   
 109 only for one pair  $(x, y)$ , and we have  $f(x', y') = 0$  for all other pairs  $(x', y')$ .

110 Indeed, one can easily check that for such morphisms  $f$ , the only morphisms  $g \leq f$  are the  
 111 morphisms which also have  $g(x', y') = 0$  for all pairs  $(x', y') \neq (x, y)$ . Such morphisms  $g$  are uniquely  
 112 described by the corresponding value  $g(x, y)$ . For every two such morphisms  $g$  and  $g'$ , depending on  
 113 whether  $g(x, y) \leq g'(x, y)$  or  $g'(x, y) \leq g(x, y)$ , we have  $g \leq g'$  or  $g' \leq g$ , i.e., the set  $\{g : g \leq f\}$  is  
 114 indeed linearly ordered.

115 Vice versa, let us prove that if a morphism has this property, then it has  $f(x, y) > 0$  only for one  
 116 pairs  $(x, y)$ . Indeed, if we have  $f(x, y) > 0$  and  $f(x', y') > 0$  for two different pairs  $(x, y) \neq (x', y')$ ,  
 117 then we would be able to construct two different morphisms  $g \leq f$  and  $g' \leq f$  for which  $g \not\leq g'$  and  
 118  $g' \not\leq g$ . Namely, we take:

- 119 •  $g(x, y) = f(x, y) > 0$  and  $g(x'', y'') = 0$  for all pairs  $(x'', y'') \neq (x, y)$ , and
- 120 •  $g'(x, y) = f(x', y') > 0$  and  $g'(x'', y'') = 0$  for all pairs  $(x'', y'') \neq (x', y')$ .

121 This contradicts our assumption that the set  $\{g : g \leq f\}$  is linearly ordered.

122 3°. Let us now describe, in ordered-category terms, morphisms  $f$  for which  $f(x, x) > 0$  for some  $a \in U$   
 123 and  $f(x', y')$  for all other pairs  $(x', y') \neq (x, x)$ .

124 Indeed, out of all morphisms described in Part 2 of this proof, such morphisms can be determined  
 125 by the additional condition that  $f \circ f = f$ . This condition is clearly satisfied for such morphisms,  
 126 while for morphisms for which  $f(x, y) > 0$  for some  $b \neq a$ , the composition  $f \circ f$  is, as one can see,  
 127 identically 0 and thus, different from  $f$ .

128 4°. One can see that two morphisms  $f$  and  $f'$  of the type described in Part 3 are connected by the  
 129 relation  $\leq$  (i.e.,  $f \leq f'$  or  $f' \leq f$ ) if and only if they correspond to the same element  $a \in U$ .

130 Thus, we can describe elements of the set  $U$  in ordered-category terms: as equivalent classes of  
 131 morphisms of the type described in Part 3 with respect to the relation  $(f \leq f') \vee (f' \leq f)$ .

132 Hence, if we have an automorphism, elements are mapped into elements in a 1-1 way, i.e., indeed  
 133 we have a bijection of the Universe of discourse.

134 5°. Let us now show that the degrees from the interval  $[0, 1]$  can also be described – modulo increasing  
 135 bijections of this interval – in ordered-category terms.

136 5.1°. Indeed, for each element  $a \in U$ , different degrees  $v \in [0, 1]$  can be associated with different  
 137 morphisms  $f$  described in Part 3 of this proof, i.e., morphisms for which:

- 138 •  $f(x, x) > 0$  for this element  $a$  and
- 139 •  $f(x', y') = 0$  for all pairs  $(x', y') \neq (x, x)$ .

140 Different degrees are then simply associated with different values  $v = f(x, x)$ .

141 This construction provides us with degrees at each element  $a \in U$ . To get a general description of  
 142 degrees, we need to relate the values corresponding to different elements  $x, x' \in U$ .

143 5.2°. Let us denote, by  $f_{x,v}$ , the morphism for which:

- 144 •  $f_{x,v}(x, x) = v$  and
- 145 •  $f_v(x', y') = 0$  for all pairs  $(x', y') \neq (x, x)$ .

We want, for every  $a \neq b$ , to connect the values  $v$  and  $w$  corresponding to functions  $f_{x,v}$  and  $f_{y,w}$ . This  
 connection comes from the following auxiliary result:

$$w \leq v \Leftrightarrow \exists f_{x \rightarrow y} \exists f_{y \rightarrow x} (f_{x \rightarrow y} \circ f_{x,v} \cdot f_{y \rightarrow x} = f_{y,w}).$$

146 Indeed, by definition of a composition, the values of the composition  $g \circ f$  cannot exceed the largest  
 147 value of each of the composed relations  $g$  and  $f$ . Thus, if  $f_{x \rightarrow y} \circ f_{x,v} \cdot f_{y \rightarrow x} = f_{y,w}$ , then the value  
 148  $f_{y,w}(b, b) = w$  cannot exceed the maximum value  $v$  of the function  $f_{x,v}$ ; thus,  $w \leq v$ .

149 Vice versa, if  $w \leq v$ , then we can take the following morphisms  $f_{x \rightarrow y}$  and  $f_{y \rightarrow x}$ :

- 150 •  $f_{x \rightarrow y}(x, y) = w$  and  $f_{x \rightarrow y}(x', y') = 0$  for all other pairs  $(x', y') \neq (x, y)$ , and, similarly,
- 151 •  $f_{y \rightarrow x}(y, x) = w$  and  $f_{y \rightarrow x}(x', y') = 0$  for all other pairs  $(x', y') \neq (y, x)$ .

152 In this case, as one can easily check, we have  $f_{x \rightarrow y} \circ f_{x,v} \cdot f_{y \rightarrow x} = f_{y,w}$ .

5.3°. Now that we know how to describe the relation  $w \leq v$  for functions  $f_{x,v}$  and  $f_{y,w}$  in ordered-category form, we can describe equality  $v = w$  between the degrees  $v$  and  $w$  corresponding to morphisms  $f_{x,v}$  and  $f_{y,w}$  as ( $v \leq w$ ) & ( $w \leq v$ ), i.e., in view of Part 5.2, as:

$$(\exists f_{x \rightarrow y} \exists f_{y \rightarrow x} (f_{x \rightarrow y} \circ f_{x,v} \cdot f_{y \rightarrow x} = f_{y,w})) \& (\exists g_{y \rightarrow x} \exists g_{x \rightarrow y} (g_{y \rightarrow x} \circ f_{y,w} \cdot g_{x \rightarrow y} = f_{x,v})).$$

153 This enables us to identify degrees  $v \in [0, 1]$  in ordered-category terms – by identifying them  
 154 with the functions  $f_{x,v}$  and taking into account the above possibility to compare degrees at different  
 155 elements  $a$ .

156 Hence, if we have an automorphism, degrees are mapped into degrees in a 1-1 and  
 157 order-preserving way, i.e., indeed we have a monotonic bijection  $H : [0, 1] \rightarrow [0, 1]$ .

158 6°. To complete the proof, we need to show how, for each morphism  $f$  and for every two elements  $a$   
 159 and  $b$ , we can describe the value  $f(x, y)$  in ordered-category terms. This will complete the proof that  
 160 the given automorphism has the form  $G_{\pi, H}$  for the mappings  $\pi$  and  $H$  as identified in Sections 4 and 5  
 161 of this proof.

6.1°. Let us first prove the following auxiliary result:

$$\exists f_{y \rightarrow x} (f_{y \rightarrow x} \circ f_{y,1} \circ f \circ f_{x,1} = f_{x,v}) \Leftrightarrow v \leq f(x, y).$$

162 Indeed, by definition of a composition, the composition  $c \stackrel{\text{def}}{=} f \circ f_{x,1}$  has the following form:

- 163 •  $c(x, y') = f(x, y')$  for all  $y'$  and
- 164 •  $c(x', y') = 0$  for all  $y'$  and for all  $x' \neq a$ .

165 Similarly, the composition  $c' \stackrel{\text{def}}{=} f_{y,1} \circ f \circ f_{x,1} = f_{y,1} \circ c$  has the following form:

- 166 •  $c'(x, y) = f(x, y)$ , and
- 167 •  $c'(x', y') = 0$  for all other pairs  $(x', y') \neq (x, y)$ .

168 As we have argued in Part 5 of this proof, the value of a composition function cannot exceed the  
 169 maximum value of each of the composed morphisms. Thus, for the composition  $f_{y \rightarrow x} \circ f_{y,1} \circ f \circ f_{x,1} =$   
 170  $f_{y \rightarrow x} \circ c'$ , the maximum value cannot exceed the maximum value  $f(x, y)$  of the morphism  $c'$ . Thus, if  
 171  $f_{y \rightarrow x} \circ c' = f_{x,v}$ , the maximum value  $v$  of the morphism  $f_{x,v}$  cannot exceed  $f(x, y)$ :  $v \leq f(x, y)$ .

172 Vice versa, for every  $v \leq f(x, y)$ , we can construct a morphism  $f_{y \rightarrow x}$  for which  $f_{y \rightarrow x} \circ c' = f_{x,v}$ :  
 173 namely, we can take:

- 174 •  $f_{y \rightarrow x}(y, x) = v$ , and
- 175 •  $f_{y \rightarrow x}(x', y') = 0$  for all pairs  $(x', y') \neq (y, x)$ .

176 One can easily check that in this case indeed  $f_{y \rightarrow x} \circ c' = f_{x,v}$ .

177 6.2°. For each morphism  $f$  and for every two elements  $a$  and  $b$ , we can identify the degree  $f(x, y)$  as  
 178 the largest degree  $v$  for which the inequality  $v \leq f(x, y)$  holds.

179 Since, according to Part 6.1 of this proof, the inequality  $v \leq f(x, y)$  can be described in  
 180 ordered-category terms, we can thus conclude that the degree  $f(x, y)$  can also be described in  
 181 ordered-category terms.

182 The proposition is proven.

#### 183 4. Conclusions

184 Many concepts of modern mathematics, starting from the basic notions of sets and functions,  
185 are described in terms of category theory. It is therefore reasonable to ask whether similar fuzzy  
186 notions can also be described in category terms. In this paper, we show that fuzzy relations – i.e., fuzzy  
187 analogues of functions – can indeed be described in category terms.

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